

Searching the Variability Space to Fix Model Inconsistencies: A Preliminary Assessment

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- ***Software Product Lines (SPL)***
 - Families of related systems distinguished by the features
 - Extensive success in achieving software reuse
 - *Variability*
 - The capacity of software artifacts to change
 - Its effective management is at the core of SPL

- ***Model-Driven Engineering (MDE)***
 - Emerging software development paradigm
 - Raises the level of abstraction and automates program generation

- **Fact**
 - Increasing research convergence in SPL and MDE that leverages their complementary capabilities

- **Challenge**
 - Maintain consistency between models with variability
 - Checking that certain relations between model elements hold for *all* products of a SPL

- **Problem**
 - Research has focused on inconsistency detection
 - Inconsistency fixing has not been fully explored

Why Search-Based SE?



■ Facts

- SPL can employ multiple models simultaneously
 - State charts, sequence diagrams, class diagrams, ...
- Models can have a large number of consistency rules and instances
- SPL usually involve large number of different products

■ Because ...

- Large search space
- No unique or optimal fixing solutions

Our ongoing work ...



- Description
 - Finds fixing locations for single constraint instances
 - Uses basic search algorithm

- Relies on the notion of Safe Composition
 - Programming languages
 - guarantee that *all* programs of a product line are type safe
 - Uses propositional logic

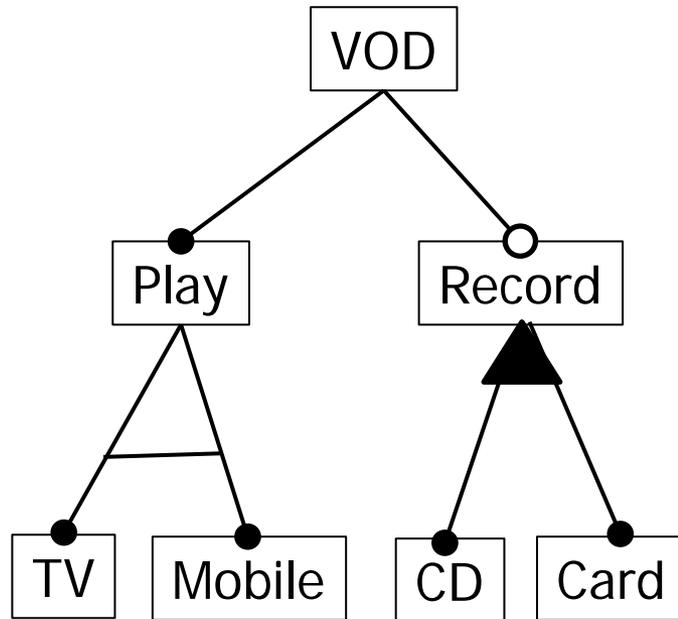
- Intuition
 - Implementation constraint(s) (IMP_f) must follow the product line domain constraints (PL_f)
- A SAT solver checks if **one** propositional formula is satisfiable or not
- Our interest is verifying that **all** the product line members satisfy an implementation constraint

$$\neg (\text{PL}_f \Rightarrow \text{IMP}_f)$$

Unsatisfiable = there is no product that violates the constraint

Satisfiable = there is at least one product that violates the constraint

Example PL_f



$VOD \Leftrightarrow \text{true} \quad \wedge$

$VOD \Leftrightarrow \text{Play} \quad \wedge$

$\text{Record} \Rightarrow VOD \quad \wedge$

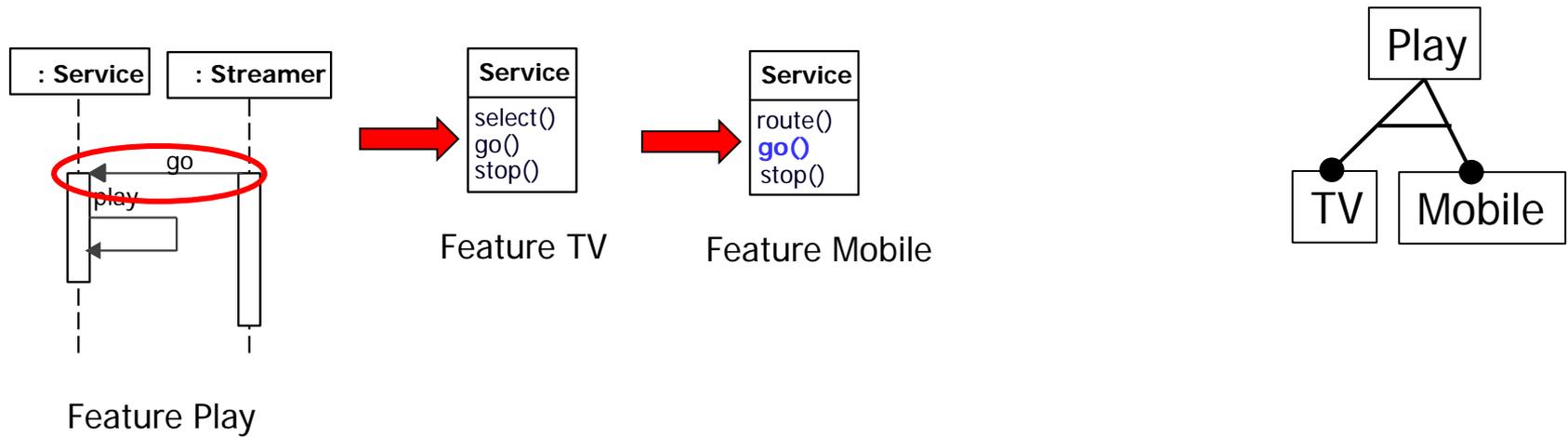
$\text{TV} \Leftrightarrow \neg \text{Mobile} \wedge \text{Play} \quad \wedge$

$\text{Mobile} \Leftrightarrow \neg \text{TV} \wedge \text{Play} \quad \wedge$

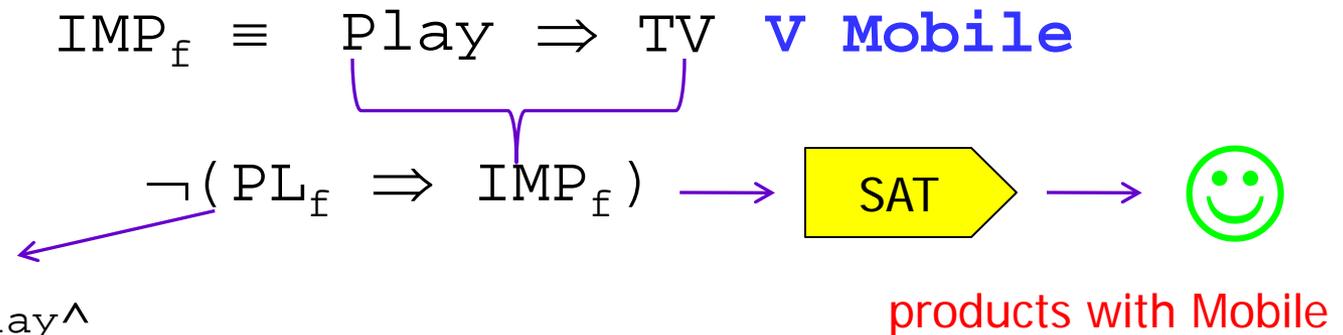
$\text{Record} \Leftrightarrow \text{CD} \vee \text{Card}$

Requiring Rule Example

- Message action must be defined as an operation in receiver's class



$VOD \Leftrightarrow true \wedge$
 $VOD \Leftrightarrow Play \wedge$
 $Record \Rightarrow VOD \wedge$
 $TV \Leftrightarrow \neg Mobile \wedge Play \wedge$
 $Mobile \Leftrightarrow \neg TV \wedge Play \wedge$
 $Record \Leftrightarrow CD \vee Card \wedge$



Fixing inconsistencies

- **Consistency Rule Instance (CRI)** is a 4-tuple $[F, RME, TS, FC]$ where:
 - Requiring feature $F = \text{Play}$
 - Requiring model element $RME = \text{play}_{\text{msg}}$
 - Set of pairs (feature, required elements) $TS = \{(TV, \text{play}_{\text{op}})\}$
 - Set of features in the pairs of TS as $TS[\text{feature}]$.
 - Faulty feature configuration that violates the consistency rule instance.
 $FC = [\{\text{VOD}, \text{Play}, \text{Mobile}\}, \{\text{TV}, \text{Record}, \text{CD}, \text{Card}\}]$
- **Pair-wise commonality**
 - Operation that receives a feature model P and two features F and G, and returns the number of products that have both features.
- **Question: where to add the required elements?**
 - Intuition: iteratively search fixing configurations choosing first those features with higher commonality → fix more configurations

Algorithm – Based on BFS

Input: CRI [F,RME,TS, FC] with $FC \neq \emptyset_{\text{conf}}$, and PL_f .

Output: Fixing set FS

$FC' := FC$

$FS := TS[\text{feature}]$

$FSQ.\text{enqueue}(FS)$

while $FC' \neq \emptyset_{\text{conf}}$ do

$FSQ.\text{dequeue}(FS)$

$G := \text{maxCom}(F, FC', TS, FS)$

$FS := FS \cup G$

$FC' := \text{SafeComposition}(PL_f, F, FS)$

$FSQ.\text{enqueue}(FS)$

end while

return FS

Set of features that guarantees no faulty configurations

Chooses a feature with maximum pair-wise commonality

Applies basic Safe Composition

- Preliminary evaluation
 - Using 60 publicly available feature models, 6-94 features
 - On average fixing sets around 5 elements
- Future work
 - Consider multiple consistency instances
 - Assess other search alternatives
 - Research alternatives for Pair-wise Commonality

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Questions?